Graphs Introduction

Definitions Structure Properties

Graphs Definition G = (V, E)

V = set of nodes

set -> unique, unordered

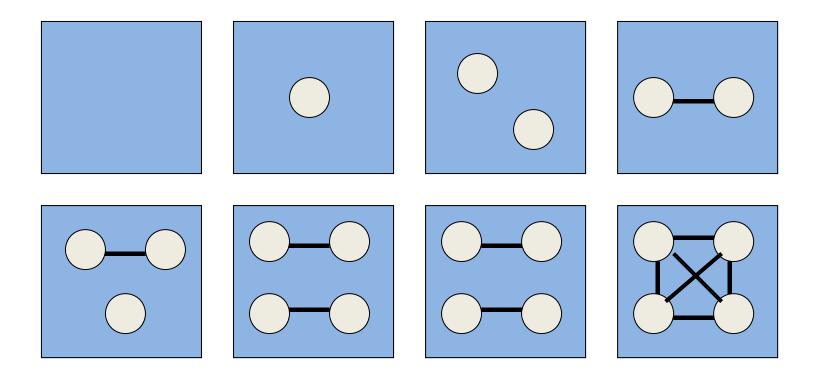
- E = set of edges (v,w) v,w nodes in V
- Degree of a node = number of incident edges
 - Directed graph in-degree / out-degree
- Edges may be
 - Undirected $v \leftarrow \rightarrow u$ undirected graph
 - Directed $v \rightarrow u$ directed graph
 - An edge connects 2 nodes represents a relationship
 - An edge may be <u>weighted</u> i.e. have a value attribute
- A graph may be empty; contain nodes only; contain nodes & edges
- Remember this definition G = (V, E)!!!

Meaning & Use

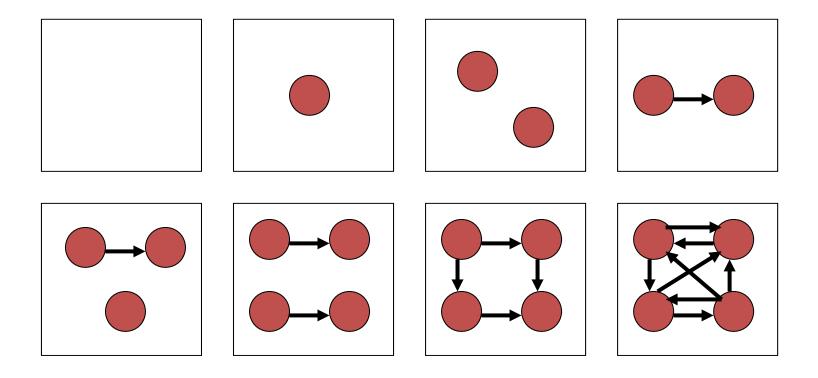
- A graph is used to represent arbitrary relationships among data objects
- e.g. undirected graphs
 - o communications network
 - o transport network (road, rail, air, sea) with costs/distances
 - o (travelling salesman problem)
- e.g. directed graphs

- (digraph)
- o flow of control in computer programs
- University course planning (dependency graph)
- state transition diagrams







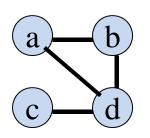


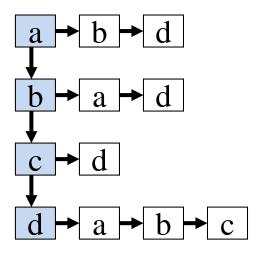
Implementations: undirected

Adjacency list

|V| + |E|

• Weights (costs may be included in the edge list)



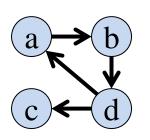


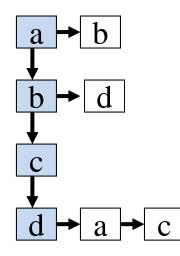


Adjacency list

|V| + |E|

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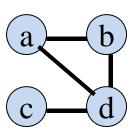


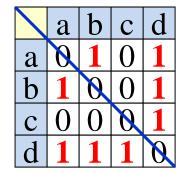


Implementations: undirected

Adjacency matrix |V| * |V|

o 4 edges - implementation 8 directed edges



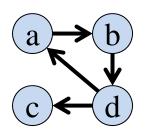


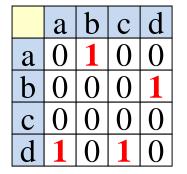
• Symmetrical about the left diagonal



Adjacency matrix |V| * |V|

o 4 edges - implementation 4 directed edges

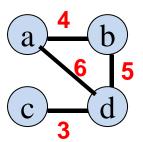


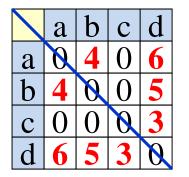


Implementations: undirected

A Cost matrix

• Weighted graphs



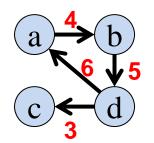


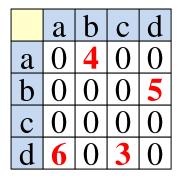
• Symmetrical about the left diagonal

Implementations: directed

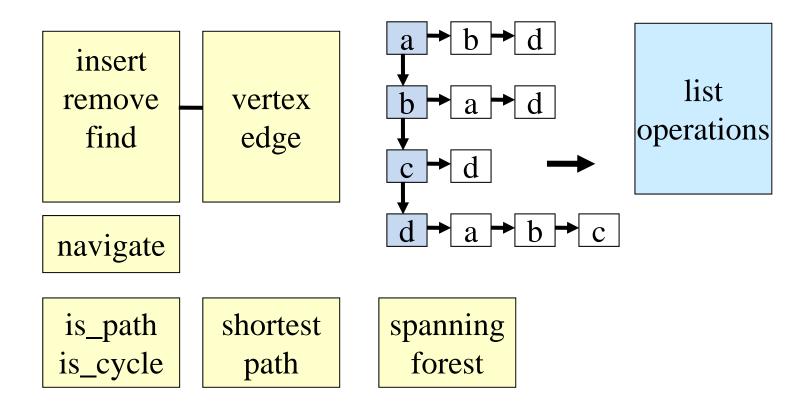
A Cost matrix

• Weighted graphs





Implementation operations



Implementation operations

- add node if node exists error else add node
- add edge (v, w) nodes v & w must exist else error
- find node found or not found
- find edge found or not found
- remove node the node must exist else error
 <u>all incident edges must also be removed</u>
- remove edge the edge must exist else error
- count nodes / count edges

Terminology

- Vertex (source / sink)
- Edge (bridge)
- Weight (positive/negative)
- Degree (in / out)
- Directed
- Undirected
- Path
- Simple path
- Cycle
- Simple cycle

- Subgraph
- Induced subgraph
- Complete graph
- Connected graph
- Connected Components
- Bipartite graph
- Spanning tree
- Spanning forest
- Shortest path tree
- DAG

Properties

- A graph with V vertices has at most V(V-1)/2 edges
- A digraph with V vertices has at most V(V-1) edges
- Path a sequence of adjacent vertices
- Simple path edges and vertices are distinct
- Simple cycle simple path where 1st/final vertices same Graph (3 vertices) / digraph (2 vertices)
- Connected graph path from every v to every other w
- An acyclic connected graph is called a TREE
 - Spanning tree / spanning forest
 - Note the difference between a directed acyclic graph (DAG) and a tree

ssues

- Finding paths (shortest A2B) and cycles
- Finding connectivity

Separability

- Edge separable (bridge edge)
- Vertex separable (articulation point)

Biconnectivity (k-edge-connected)

- Every pair of vertices connected by 2 disjoint paths 0
- No articulation points Ο

Algorithms

- Is_path(a, b)
- Is_cycle(G)
- Is_connected(G)
- Is_strongly_connected(DG)
- Transitive closure (G)
- Minimal spanning tree (MST) Prim / Kruskal
- Single-source shortest paths
- All-pairs shortest paths
- Topological sort (DAG)
- Depth / breadth first search

Directed Graph Warshall

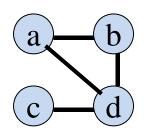
Dijkstra

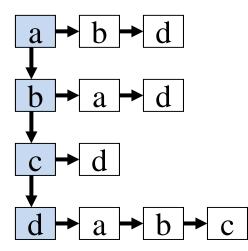
- Floyd
- **Directed Acyclic Graph**

Traversing a graph

- Mark each node when visited
- **Depth-first** search: $a \rightarrow b \rightarrow d \rightarrow c$

Breadth-first search: a, one step (b, d), two steps C





Visiting Nodes – 2 methods

- 1. Mark as visited (sometimes recursive calls)
- 2. **Cut**
 - Divide the graph into <u>components</u> (n nodes)
 - Merge the components by adding edges
 - Independent components (Kruskal) (using a PQ)
 - Choose one **start node** (in a component)
 - <u>2 sets</u>: visited (S) and not visited (V-S)
 - Merge the components by adding edges
 Tree formation (Prim, Dijkstra)

Depth-first search (dfs)

Recursive search

(stack)

- Depth-first numbering
 - o Preorder
 - o Postorder
- Cost
 - O(|V|+|E|)
 - O(V²)

(order that processing starts) (order that processing finishes)

(adjacency list) (adjacency matrix)

Depth-first search (dfs) used for

- Simple path
- Simple connectivity (dfs called once)
- Topological sort (digraphs DAGs)
- Finding strongly connected components
- Cycle detection (back edges)
- Finding bridge edges
- Finding articulation points

Digraph (Directed Graph) Algorithms

- dfs / bfs O(|V|+|E|) $O(V^2)$
- $O(V^2)$ Dijkstra
- $O(V^3)$ Floyd

 $O(V^3)$ Warshall

(cheaper variants exist)

- (cheaper variants exist)
- Alternative dfs on each node $O(V^2)$
- Topological sort (DAG) dfs O(|V|+|E|)
- Strong components (dfs)
 - Kosaraju O(|V|+|E|) or $O(V^2)$ 0
 - Tarjan 0
 - Gabow (1999) \mathbf{O} (we will not consider these 3 algorithms in this course)

Graph Algorithms

- dfs / bfs $O(|V|+|E|) O(V^2)$
- MST
 - **Prim** O(V²) (1961)
 - Kruskal O(E Ig E) (1956)
- Other problems
 - Travelling Salesman (Hamiltonian path)
 - Königsberg Bridges (Euler)
 - Matching (bipartite graphs)

Summary

- Collection of nodes + relationships
- G = (V,E) directed; undirected
- Edges may be weighted
- Implementation: adjacency list, matrix
- Digraphs: represent dependencies
- Graphs: represent networks