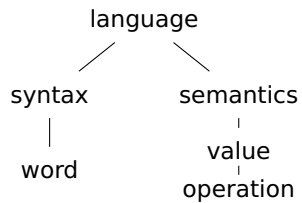


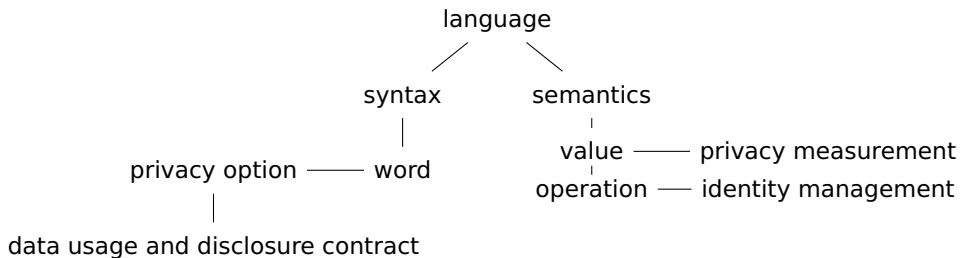
Towards a Formal Language for Privacy Options

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$\underbrace{c_1 \text{ 'predicate' } c_2}_{\text{Contract}}$ } combination of contracts

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e. g., postponing of contract execution

$\underbrace{c_1 \text{ 'predicate' } c_2}_{\text{Contract}}$ } combination of contracts
e. g., Boolean operators

$c_1 = \text{zero}$

... no rights, no obligations

$c_2 = \text{data } a_1 p_1$

... receive data and use it immediately

$c_3 = \text{give } c_2$

... send data which is used immediately

$c_4 = c_1 \text{ 'or' } c_2$

... receive and use the data or do nothing

$c_5 = \text{when (at } t) c_4$

... postpone c_4 to time t

zero :: Contract

data :: Contract

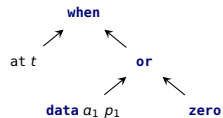
give :: Contract



and :: Contract

or :: Contract

cond :: Contract



when :: Contract

anytime :: Contract

until :: Contract



$$ACTION = \text{Date} \rightarrow IO \text{ Contract}$$

$$\mathcal{O}_{\mathcal{E}}[\cdot] : \text{Contract} \rightarrow ACTION$$

$$\mathcal{O}_{\mathcal{E}}[\text{zero}] = \text{return}(\text{zero})$$

$$\mathcal{O}_{\mathcal{E}}[\text{data } a \ p] = \text{use}(a, p)$$

$$\mathcal{O}_{\mathcal{E}}[\text{give } c] = \text{send}(c)$$

$$\mathcal{O}_{\mathcal{E}}[c_1 \ \text{'and'} \ c_2] = \mathcal{O}_{\mathcal{E}}[c_1] \parallel \mathcal{O}_{\mathcal{E}}[c_2]$$

$$\mathcal{O}_{\mathcal{E}}[c_1 \ \text{'or'} \ c_2] = \text{greedy}_{\mathcal{E}}(\mathcal{O}_{\mathcal{E}}[c_2], \mathcal{O}_{\mathcal{E}}[c_1])$$

$$\mathcal{O}_{\mathcal{E}}[\text{cond } o \ c_1 \ c_2] = \text{if}(\mathcal{V}[o], \mathcal{O}_{\mathcal{E}}[c_1], \mathcal{O}_{\mathcal{E}}[c_2])$$

$$\mathcal{O}_{\mathcal{E}}[\text{when } o \ c] = \text{when}'_{\mathcal{E}}(\mathcal{V}[o], \mathcal{O}_{\mathcal{E}}[c])$$

$$\mathcal{O}_{\mathcal{E}}[\text{anytime } o \ c] = \text{stopping}'_{\mathcal{E}}(\mathcal{V}[o], \mathcal{O}_{\mathcal{E}}[c])$$

$$\mathcal{O}_{\mathcal{E}}[\text{until } o \ c] = \text{absorb}'_{\mathcal{E}}(\mathcal{V}[o], \mathcal{O}_{\mathcal{E}}[c])$$

Motivation

- language with similar design goals: PPL Sticky Policies.
- $POL \not\subseteq PPL$ and $POL \not\supseteq PPL$, thus no easy comparison.
- way out: $POL^- \subseteq POL$ with $POL^- \subseteq PPL$.

Q: how to define POL^- ?

- when and anytime *not* part of POL^- .
- points of time the *only* conditions of `until` and `cond`.
- new PPL action: 'immediate usage'.

- add new language primitives (experts only).
 - experts only.
 - risk for redundancy.
 - readjustment of semantics.
- combinators for common contract patterns (common user).
 - common user.
 - application of existing language primitives.
 - no readjustment of semantics.
 - ex def: `twice f c = (f c) 'and' (f c)`.
 - ex app: `until (at t) 'twice' c1`

- def: language for describing privacy options (syntax, semantics).
- word in that language can be interpreted as:
 - (numeric) value – privacy measurement.
 - sequence of i/o action – identity management.
- comparison to existing languages (PPL).
- discussion: extending the language.

Q & A

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Recommended reading. Berthold, S., Böhme, R.: Valuating privacy with option pricing theory. In: Workshop on the Economics of Information Security (WEIS), University College London, UK (2009).